### Contextuality as a resource

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Combining Viewpoints in Quantum Theory Edinburgh, 20th March 2018 Joint work with:

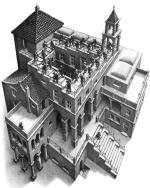
- Samson Abramsky (Oxford)
- Shane Mansfield (Paris VII)

and also:

- Kohei Kishida (Dalhousie)
- Giovanni Carù (Oxford)
- Nadish de Silva (UCL)
- Octavio Zapata (UCL)

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fundamental non-classical phenomenona of QM



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### MSD

Howard-Wallman-Veith-Emerson (2014)

"Contextuality supplies the 'magic' for quantum computation"

Contextuality formulated in a theory-independent fashion

 Abramsky & Brandenburger: unified framework for non-locality and contextuality (cf. Cabello–Severini–Winter, Acín–Fritz–Leverrier–Sainz)

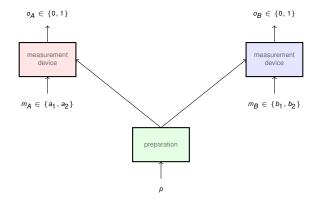
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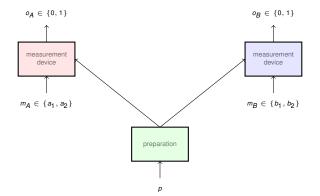
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  - Quantifiable advantages in QC and QIP tasks

# **Empirical data**



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Using elementary logic and probability:

$$1 = \operatorname{Prob}(\neg \bigwedge \phi_i) = \operatorname{Prob}(\bigvee \neg \phi_i)$$
$$\leq \sum_{i=1}^{N} \operatorname{Prob}(\neg \phi_i) = \sum_{i=1}^{N} (1 - p_i) = N - \sum_{i=1}^{N} p_i.$$

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• Hence,  $\sum_{i=1}^{N} p_i \le N - 1$ .

			( <mark>0</mark> , 1)		
		1/2	0	0	1/2
$a_1$	<b>b</b> <sub>2</sub>		1/8	1/8	<sup>3</sup> /8
$a_2$	$b_1$	<sup>3</sup> /8	1/8		
<b>a</b> 2	<b>b</b> 2	1/8	3/8	<sup>3</sup> /8	1/8

Α	В	( <mark>0,0</mark> )	( <mark>0</mark> , 1)	(1, <mark>0</mark> )	<b>(1</b> , <b>1</b> )
$a_1$	$b_1$	1/2	0	0	1/2
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$a_2$	$b_1$	3/8	1/8	1/8	<sup>3</sup> /8
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These formulae are contradictory.

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The inequality is violated by 1/4.

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- What was our unwarranted assumption?
- That all variables could in principle be observed simultaneously.
- Local consistency vs global inconsistency.

## Abramsky–Brandenburger framework

Measurement scenario  $\langle X, \mathcal{M}, O \rangle$ :

- X is a finite set of measurements or variables
- O is a finite set of outcomes or values
- ▶ *M* is a cover of *X*, indicating **joint measurability** (contexts)

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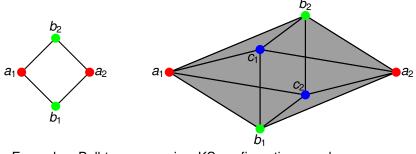
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Example: (2,2,2) Bell scenario

- The set of variables is  $X = \{a_1, a_2, b_1, b_2\}$ .
- The outcomes are  $O = \{0, 1\}$ .
- The measurement contexts are:

 $\{ \{a_1, b_1\}, \ \{a_1, b_2\}, \ \{a_2, b_1\}, \ \{a_2, b_2\} \}.$ 

# Measurement scenarios



Examples: Bell-type scenarios, KS configurations, and more.

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- A set of 18 variables, X = {A,..., O}
- A set of outcomes *O* = {0, 1}
- ► A measurement cover *M* = {*C*<sub>1</sub>,..., *C*<sub>9</sub>}, whose contexts *C<sub>i</sub>* correspond to the columns in the following table:

$U_1$	$U_2$	$U_3$	$U_4$	$U_5$	$U_6$	$U_7$	$U_8$	U <sub>9</sub>
Α	Α	Н	Н	В	1	Р	Р	Q
В	Е	1	K	Е	Κ	Q	R	R
С	F	С	G	М	Ν	D	F	М
D	G	J	L	Ν	0	J	L	0

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**Empirical model**: family  $\{e_C\}_{C \in \mathcal{M}}$  where  $e_C \in \operatorname{Prob}(O^C)$  for  $C \in \mathcal{M}$ .

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Compatibility condition: the distributions "agree on overlaps"

$$\forall C, C' \in \mathcal{M}. \quad e_C|_{C \cap C'} = e_{C'}|_{C \cap C'}.$$

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In multipartite scenarios, compatibility = the **no-signalling** principle.

A (compatible) empirical model is **non-contextual** if there exists a **global distribution**  $d \in \text{Prob}(O^X)$  on the joint assignments of outcomes to all measurements that marginalises to all the  $e_C$ :

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The import of results such as Bell's and Bell–Kochen–Specker's theorems is that there are empirical models arising from quantum mechanics that are contextual.

#### Possibilistic collapse

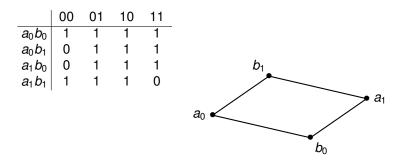
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- Contains the possibilistic, or logical, information of that model.

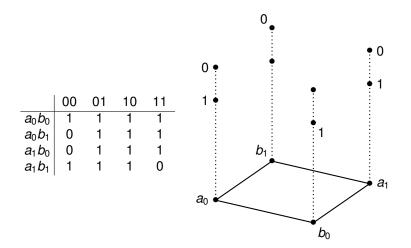
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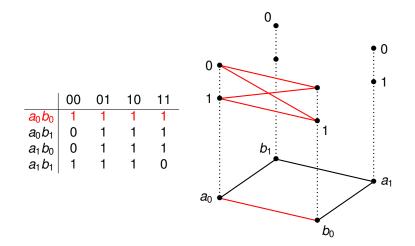
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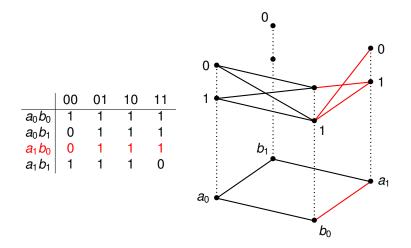
	00	01	10	11			00	01	10	11
$a_1b_1$	1/2	0	0	1/2		$a_1b_1$	1	0	0	1
$a_1b_2$	3/8	1/8	1/8	3/8	$\mapsto$	$a_1b_2$	1	1	1	1
$a_2b_1$	3/8	1/8	1/8	3/8		$a_2b_1$	1	1	1	1
$a_2b_2$	1/8	3/8	3/8	1/8		$a_2b_2$	1	1	1	1

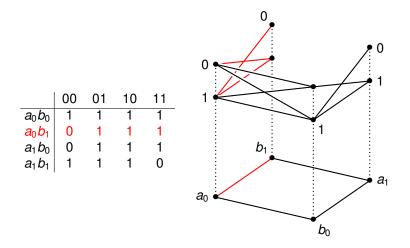
	00	01	10	11	
$a_0b_0$	1	1	1	1	
$a_0b_1$	0	1	1	1	
$ \begin{array}{c} a_0 b_0 \\ a_0 b_1 \\ a_1 b_0 \end{array} $	0	1	1	1	
$a_1b_1$	1	1	1	0	

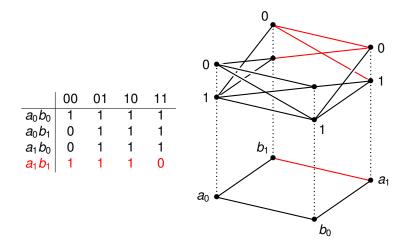


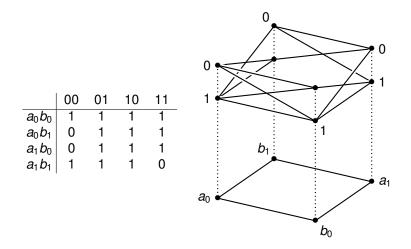


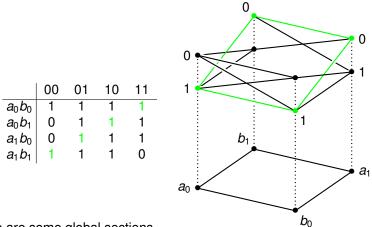




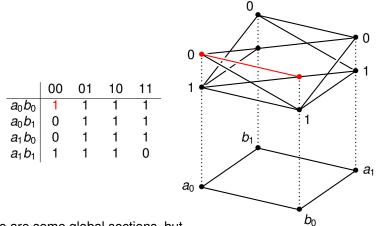


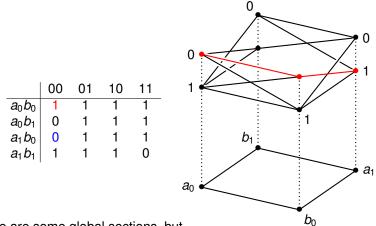


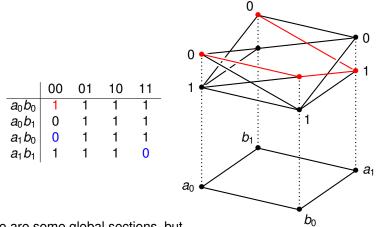


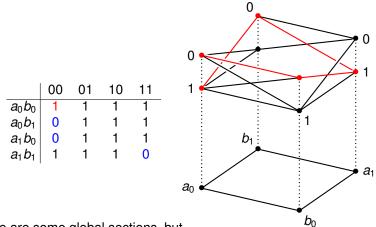


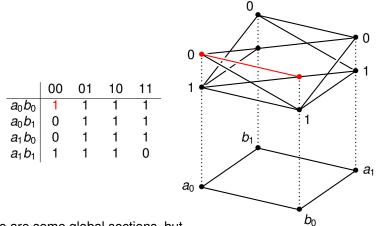
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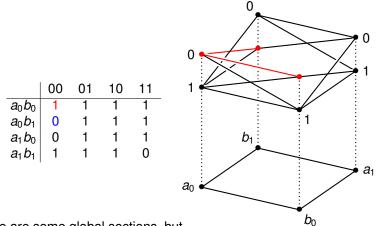


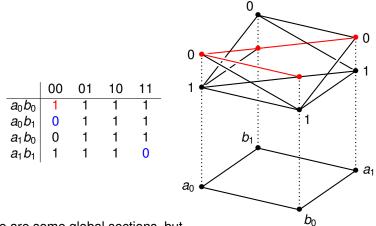


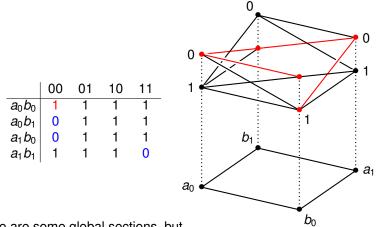


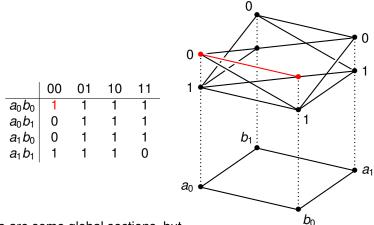












There are some global sections, but ...

Logical contextuality: Not all sections extend to global ones.

#### Strong contextuality

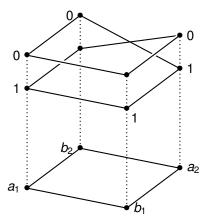
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E.g. K-S, GHZ, the PR box:

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$a_1$	$b_1$	1	0	0	1
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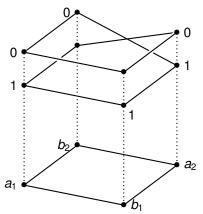


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Cohomological witnesses of contextuality (Abramsky–B–Mansfield, ABM–Kishida–Lal, Carù, Raussendorf et al.)

# Measuring Contextuality

Non-contextuality: global distribution  $d \in \text{Prob}(O^X)$  such that:

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Consider **subdistributions**  $c \in \text{SubProb}(O^{\chi})$  such that:

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Equivalently, maximum weight  $\lambda$  over all convex decompositions

$$m{e} = \lambda m{e}^{NC} + (1 - \lambda)m{e}'$$

where  $e^{NC}$  is a non-contextual model.

## The contextual fraction

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$$\boldsymbol{e} = \lambda \boldsymbol{e}^{\boldsymbol{N}\boldsymbol{C}} + (1-\lambda)\boldsymbol{e}^{\boldsymbol{S}\boldsymbol{C}}$$

where  $e^{NC}$  is a non-contextual model.  $e^{SC}$  is strongly contextual!

$$NCF(e) = \lambda$$
  $CF(e) = 1 - \lambda$ 

# (Non-)contextual fraction via linear programming

Checking contextuality of e corresponds to solving

Find 
$$\mathbf{d} \in \mathbb{R}^n$$
  
such that  $\mathbf{M}\mathbf{d} = \mathbf{v}^e$   
and  $\mathbf{d} \ge \mathbf{0}$ .

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Computing the non-contextual fraction corresponds to solving the following linear program:

Find	$\mathbf{c} \in \mathbb{R}^n$
maximising	1 · c
subject to	$\mathbf{Mc}\leq\mathbf{v}^{e}$
and	$\mathbf{c} \geq 0$

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# E.g. Equatorial measurements on GHZ(n)

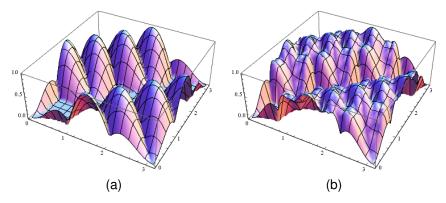


Figure: Contextual fraction of empirical models obtained with equatorial measurements at  $\phi_1$  and  $\phi_2$  on each qubit of  $|\psi_{\text{GHZ}(n)}\rangle$  with: (a) n = 3; (b) n = 4.

# Violations of Bell inequalities

An **inequality** for a scenario  $\langle X, \mathcal{M}, O \rangle$  is given by:

- A set of coefficients α = {α(C, s)}<sub>C∈M,s∈O<sup>C</sup></sub>
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For a model *e*, the inequality reads as

$$\mathcal{B}_{lpha}(oldsymbol{e})\ \leq\ oldsymbol{R}$$
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where

$$\mathcal{B}_{lpha}(m{e}) \ \coloneqq \ \sum_{m{C}\in\mathcal{M},m{s}\in O^{\mathcal{C}}} lpha(m{C},m{s})m{e}_{m{C}}(m{s}) \ .$$

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NB: A complete set of inequalities can be derived from the logical approach.

# Violation of a Bell inequality

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ight\}$$

The **normalised violation** of a Bell inequality  $\langle \alpha, R \rangle$  by an empirical model *e* is the value

$$rac{\max\{\mathbf{0},\mathcal{B}_lpha(m{e})-m{R}\}}{\|lpha\|-m{R}}\;.$$

#### Proposition

Let *e* be an empirical model.

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- ► The normalised violation by *e* of any Bell inequality is at most CF(*e*).
- This bound is attained: there exists a Bell inequality whose normalised violation by e is exactly CF(e).
- Moreover, this Bell inequality is tight at "the" non-contextual model e<sup>NC</sup> and maximally violated by "the" strongly contextual model e<sup>SC</sup> for any decomposition:

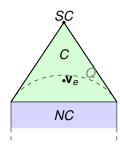
$$e = \mathsf{NCF}(e)e^{\mathsf{NC}} + \mathsf{CF}(e)e^{\mathsf{SC}}$$

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#### Quantifying Contextuality LP:

$\mathbf{c} \in \mathbb{R}^n$
1 · c
${f M}{f c}\leq{f v}^{e}$
$\textbf{c} \geq \textbf{0}$

$$\boldsymbol{e} = \lambda \boldsymbol{e}^{NC} + (1 - \lambda) \boldsymbol{e}^{SC}$$
 with  $\lambda = \mathbf{1} \cdot \mathbf{x}^*$ .



#### Quantifying Contextuality LP:

Find  $\mathbf{C} \in \mathbb{R}^n$  maximising  $\mathbf{1} \cdot \mathbf{C}$ 

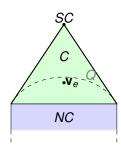
subject to  $MC \leq v^e$ 

and  $\mathbf{c} \geq \mathbf{0}$ 

 $\boldsymbol{e} = \lambda \boldsymbol{e}^{NC} + (1 - \lambda) \boldsymbol{e}^{SC}$  with  $\lambda = \mathbf{1} \cdot \mathbf{x}^*$ .

Dual LP:

Find	$\mathbf{y} \in \mathbb{R}^m$
minimising	y · v <sup>e</sup>
subject to	$\mathbf{M}^{T}\mathbf{y} \geq 1$
and	$\mathbf{y} \geq 0$



#### Quantifying Contextuality LP:

Find  $\mathbf{c} \in \mathbb{R}^n$ maximising  $\mathbf{1} \cdot \mathbf{c}$ subject to  $\mathbf{M} \mathbf{c} \leq \mathbf{v}^e$ 

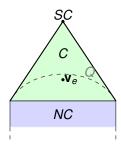
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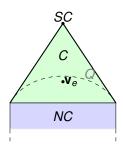
$$\textbf{a} \mathrel{\mathop:}= \textbf{1} - |\mathcal{M}|\textbf{y}|$$



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Find	$\mathbf{a} \in \mathbb{R}^m$
maximising	a · v <sup>e</sup>
subject to	M <sup>7</sup> a ≤ <b>0</b>
and	a ≤ 1

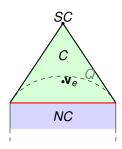
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# $\textbf{a} \mathrel{\mathop:}= \textbf{1} - |\mathcal{M}|\textbf{y}$

Find	$\mathbf{a} \in \mathbb{R}^m$
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subject to	M <sup>7</sup> a ≤ <b>0</b>
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.

computes tight Bell inequality (separating hyperplane)

# Operations on empirical models

## Contextuality as a resource

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- More than one possible measure of contextuality.
- What properties should a good measure satisfy?

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- More than one possible measure of contextuality.
- What properties should a good measure satisfy?

- Monotonicity wrt operations that do not introduce contextuality
- Towards a resource theory as for entanglement (e.g. LOCC), non-locality, ...

Think of empirical models as black boxes

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What operations can we perform (non-contextually) on them?

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- We write type statements

 $e:\langle X,\mathcal{M},\mathcal{O}
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to mean that *e* is a (compatible) emprical model on  $\langle X, \mathcal{M}, O \rangle$ .

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to mean that *e* is a (compatible) emprical model on  $\langle X, \mathcal{M}, O \rangle$ .

The operations remind one of process algebras.



Relabelling

$$\begin{array}{ll} \boldsymbol{e}:\langle \boldsymbol{X},\mathcal{M},\boldsymbol{O}\rangle\\ \alpha:(\boldsymbol{X},\mathcal{M})\cong(\boldsymbol{X}',\boldsymbol{M}') & \rightsquigarrow \boldsymbol{e}[\alpha]:\langle \boldsymbol{X}',\mathcal{M}',\boldsymbol{O}\rangle\end{array}$$

Relabelling

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For 
$$C \in \mathcal{M}, s : \alpha(C) \longrightarrow O, e[\alpha]_{\alpha(C)}(s) := e_C(s \circ \alpha^{-1})$$

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Restriction

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#### Restriction

$$\begin{array}{ll} \boldsymbol{e}:\langle \boldsymbol{X}, \mathcal{M}, \boldsymbol{O} \rangle \\ (\boldsymbol{X}', \mathcal{M}') \leq (\boldsymbol{X}, \boldsymbol{M}) \end{array} \rightsquigarrow \boldsymbol{e} \upharpoonright \mathcal{M}': \langle \boldsymbol{X}', \mathcal{M}', \boldsymbol{O} \rangle \end{array}$$

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Coarse-graining

$$\begin{array}{ll} e: \langle X, \mathcal{M}, O \rangle \\ f: O \longrightarrow O' & \rightsquigarrow e/f: \langle X, \mathcal{M}, O' \rangle \end{array}$$

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For 
$$C \in M, s : C \longrightarrow O', (e/f)_C(s) := \sum_{t:C \longrightarrow O, f \circ t = s} e_C(t)$$

$$\begin{array}{ll} \text{Mixing} & \begin{array}{c} \boldsymbol{e}, \boldsymbol{e}' : \langle \boldsymbol{X}, \mathcal{M}, \boldsymbol{O} \rangle \\ \lambda \in [0, 1] \end{array} \quad \rightsquigarrow \quad \boldsymbol{e} +_{\lambda} \boldsymbol{e}' : \langle \boldsymbol{X}, \mathcal{M}, \boldsymbol{O} \rangle \end{array}$$

# Operations Mixing

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$$\begin{array}{l} \mathsf{For}\; {\pmb{C}} \in {\pmb{M}}, {\pmb{s}}: {\pmb{C}} \longrightarrow {\pmb{O}}', \\ ({\pmb{e}}+_{\lambda} \; {\pmb{e}}')_{{\pmb{C}}}({\pmb{s}}) := \lambda {\pmb{e}}_{{\pmb{C}}}({\pmb{s}}) + (1-\lambda) {\pmb{e}}'_{{\pmb{C}}}({\pmb{s}}) \end{array}$$

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Choice

$$egin{array}{lll} e:\langle X,\mathcal{M},\mathcal{O}
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Choice

$$\begin{array}{ll} e: \langle X, \mathcal{M}, O \rangle \\ e': \langle X', \mathcal{M}', O \rangle \end{array} \rightsquigarrow e \& e': \langle X \sqcup X', \mathcal{M} \sqcup \mathcal{M}', O \rangle \end{array}$$

For  $C \in M$ ,  $(e \& e')_C := e_C$ For  $D \in M'$ ,  $(e \& e')_D := e'_D$ 

Mixing

$$e, e' : \langle X, \mathcal{M}, O \rangle$$
  $\rightsquigarrow$   $e +_{\lambda} e' : \langle X, \mathcal{M}, O \rangle$   
 $\lambda \in [0, 1]$ 

$$\begin{array}{l} \mathsf{For}\; C \in \mathit{M}, \mathit{s}: C \longrightarrow O', \\ (\mathit{e}_{+\lambda} \; \mathit{e}')_{\mathit{C}}(\mathit{s}) := \lambda \mathit{e}_{\mathit{C}}(\mathit{s}) + (1 - \lambda) \mathit{e}_{\mathit{C}}'(\mathit{s}) \end{array}$$

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Tensor

$$\begin{array}{ll} \boldsymbol{e}:\langle \boldsymbol{X}, \mathcal{M}, \boldsymbol{O} \rangle \\ \boldsymbol{e}':\langle \boldsymbol{X}', \mathcal{M}', \boldsymbol{O} \rangle \end{array} \quad \rightsquigarrow \quad \boldsymbol{e} \otimes \boldsymbol{e}':\langle \boldsymbol{X} \sqcup \boldsymbol{X}', \mathcal{M} \star \mathcal{M}', \boldsymbol{O} \rangle \end{array}$$

Mixing

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Tensor

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$$\mathcal{M}\star\mathcal{M}':=\{\mathcal{C}\sqcup\mathcal{D}\mid\mathcal{C}\in\mathcal{M},\mathcal{D}\in\mathcal{M}'\}$$

Mixing

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$$\begin{split} \mathcal{M} \star \mathcal{M}' &:= \{ C \sqcup D \mid C \in \mathcal{M}, D \in \mathcal{M}' \} \\ \mathsf{For} \; C \in \mathcal{M}, D \in \mathcal{M}', s = \langle s_1, s_2 \rangle : C \sqcup D \longrightarrow O, \\ (e \otimes e')_{C \sqcup D} \langle s_1, s_2 \rangle &:= e_C(s_1) \, e'_D(s_2) \end{split}$$

Relabelling  $e[\alpha]$ 

Relabelling  $e[\alpha]$ 

**Restriction**  $e \upharpoonright \mathcal{M}'$ 

Relabelling  $e[\alpha]$ 

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Coarse-graining e/f

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Choice e & e'

R S Barbosa Contextuality as a resource 27

Relabelling  $e[\alpha]$ 

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Mixing $\lambda e + (1 - \lambda)e'$ Choicee & e'

Tensor  $e_1 \otimes e_2$ 

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Sequencing

Relabelling  $CF(e[\alpha]) = CF(e)$ 

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Tensor  $e_1 \otimes e_2$ 

Relabelling  $CF(e[\alpha]) = CF(e)$ 

**Restriction**  $CF(e \upharpoonright M') \leq CF(e)$ 

Coarse-graining e/f

Mixing  $\lambda \boldsymbol{e} + (1-\lambda)\boldsymbol{e}'$ 

Choice *e* & *e*'

Tensor  $e_1 \otimes e_2$ 

Relabelling  $CF(e[\alpha]) = CF(e)$ 

**Restriction**  $CF(e \upharpoonright M') \leq CF(e)$ 

Coarse-graining  $CF(e/f) \leq CF(e)$ 

Mixing  $\lambda e + (1 - \lambda)e'$ 

Choice e & e'

Tensor  $e_1 \otimes e_2$ 

 $CF(e[\alpha]) = CF(e)$ Relabelling

 $\mathsf{CF}(e \upharpoonright \mathcal{M}') \leq \mathsf{CF}(e)$ Restriction

Coarse-graining CF(e/f) < CF(e)

Mixing  $CF(\lambda e + (1 - \lambda)e') < \lambda CF(e) + (1 - \lambda)CF(e')$ 

Choice e&e'

Tensor  $e_1 \otimes e_2$ 

Sequencing

e1; e2

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## Resource theory of contextuality

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# Contextual fraction and quantum advantages

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We have

$$1-\bar{p}_S \geq \text{NCF} \frac{n-k}{n}$$

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- Probabilistic version: non-linear function computed with sufficently large probability of success implies contextuality.

▶ **Goal**: Compute Boolean function  $f : 2^m \longrightarrow 2^l$  using l2-MBQC

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where for Boolean functions *f* and *g*,  $d(f, g) := 2^{-m} | \{i \in 2^m | f(i) \neq g(i)\}$ .

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▶ Then,

$$1 - \bar{p}_S \geq \operatorname{NCF}(e) \nu(f)$$

#### Questions...



"The contextual fraction as a measure of contextuality" Samson Abramsky, RSB, Shane Mansfield PRL 119:050504 (2017), arXiv:1705.07918[quant-ph]