

Effectus theory

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Teaser, another reconstruction of QT:

an Operational f -effectus is (a subcategory of) Euclidean Jordan Algebras (with positive maps)

(finite-dimensional)

follows from [Wetering, arXiv:1801.05798]

People



Jacobs



Cho



A. Westerbaan

(≠)



B. W.

People



Jacobs



Cho



A. Westerbaan

(\neq)



B. W.



Tull
connection with
OPTs



Wetering
Reconstruction



Adams
Type theory

Oxford CQM

f.d. Hilbert spaces

\mathbb{C}^2 (qubit), \mathbb{C}^3 (qutrit), ...

Effectus theory

von Neumann algebras

\mathbb{C}^2 (bit), M_2 (qubit), ...

Oxford CQM

f.d. Hilbert spaces

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Operators

\mathbb{H} , \mathcal{A} , ...

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von Neumann algebras

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contractive normal c.p. maps

$a \mapsto \sum_i b_i^* a b_i$, ...

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\otimes parallel composition

$\mathbb{C}^2 \otimes \mathbb{C}^2$: two qubits

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$+$ probabilistic disjunction

$\mathbb{C}^2 \oplus M_2$: bit or qubit

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Operators

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$\mathbb{C}^2 \otimes \mathbb{C}^2$: two qubits

Expressive calculus for 'circuits'

Works best finite-dimensionally

Effectus theory

von Neumann algebras

\mathbb{C}^2 (bit), M_2 (qubit), ...

contractive normal c.p. maps

$a \mapsto \sum_i b_i^* a b_i$, ...

$+$ probabilistic disjunction

$\mathbb{C}^2 \oplus M_2$: bit or qubit

Hard to reason about circuits

Measurement, classical data
and infinite dimensions built in.

Maps between (vN) algebras go the opposite way

measure in
std. basis

$$\mathbb{C}^2 \rightarrow M_2$$

(bit ← qbit)

$$(\lambda, \mu) \mapsto \lambda|0\rangle\langle 0| + \mu|1\rangle\langle 1|$$

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$$M_2 \rightarrow \mathbb{C}$$

(qbit \leftarrow 1)

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$$M_2 \rightarrow M_2 \otimes M_3$$

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$$a \mapsto a \otimes 1$$

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$$a \mapsto a \otimes 1$$

CPTP

$$B(\mathcal{Y}) \rightarrow B(\mathcal{X})$$

$$\rho \mapsto \sum_i b_i^* \rho b_i$$

cpu-map

$$B(\mathcal{K}) \rightarrow B(\mathcal{Y})$$

$$a \mapsto \sum_i b_i a b_i^*$$

Overview

operational
t-effectus

is a

t-effectus

Overview

operational
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is an

effectus

An Effectus has

. objects A, B, X, Y, \dots representing data types/systems

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- final object 1 representing the system with one state

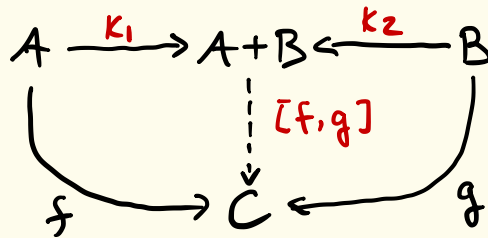
for any A , there is a unique $!_A: A \dashrightarrow 1$

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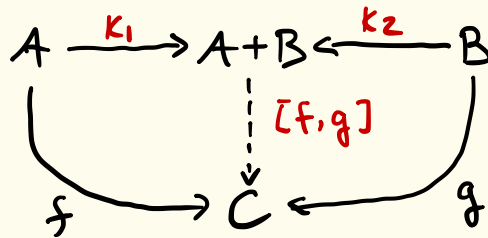


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$$\text{swap} \equiv [k_2, k_1]: A+A \rightarrow A+A$$

$$\vee \equiv [\text{id}, \text{id}]: A+A \rightarrow A$$

Predicates and partial maps

$$p: X \rightarrow \{+1\}$$

predicate on X

Predicates and partial maps

$$p: X \rightarrow 1+1$$

$$\begin{array}{ccc} X & \xrightarrow{!} & 1 \\ & \searrow & \nearrow \scriptstyle K_1 \\ & \xrightarrow{p} & 1+1 \end{array}$$

predicate on X

the truth predicate on X

Predicates and partial maps

$$p: X \rightarrow 1+1$$

$$X \xrightarrow{!} 1 \xrightarrow{K_1} 1+1$$

$\xrightarrow{\quad 1 \quad}$

$$X \xrightarrow{p} 1+1 \xrightarrow{\text{swap}} 1+1$$

$\xrightarrow{\quad p^\perp \quad}$

predicate on X

the truth predicate on X

the negation/orthosupplement of p

Predicates and partial maps

$$p: X \rightarrow 1+1$$

$$X \begin{array}{c} \xrightarrow{!} 1 \\ \xrightarrow{\quad} 1+1 \end{array} \begin{array}{c} \xrightarrow{K_1} \\ \xrightarrow{\quad} \end{array} 1+1$$

$$X \begin{array}{c} \xrightarrow{p} 1+1 \\ \xrightarrow{\quad} 1+1 \end{array} \begin{array}{c} \xrightarrow{\text{swap}} \\ \xrightarrow{\quad} \end{array} 1+1$$

\perp
 p^\perp

$$X \xrightarrow{f} Y+1$$

predicate on X

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the negation/orthosupplement of p

a partial map $X \rightarrow Y$

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p^\perp

predicate on X

the truth predicate on X

the negation/orthosupplement of p

$$X \xrightarrow{f} Y+1$$

$$X \begin{array}{c} \xrightarrow{f} Y+1 \\ \xrightarrow{\quad} Z+1 \end{array} \xrightarrow{[g, k_2]} Z+1$$

$g \circ f$

a partial map $X \rightarrow Y$

composition of partial maps

Predicates and partial maps

$$p: X \rightarrow I+1$$

$$X \begin{array}{c} \xrightarrow{i} I \xrightarrow{k_1} I+1 \\ \xrightarrow{\quad} I+1 \end{array}$$

$$X \begin{array}{c} \xrightarrow{p} I+1 \xrightarrow{\text{swap}} I+1 \\ \xrightarrow{\quad} I+1 \end{array}$$

predicate on X

the truth predicate on X

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$$X \xrightarrow{f} Y+1$$

$$X \begin{array}{c} \xrightarrow{f} Y+1 \xrightarrow{[g, k_2]} Z+1 \\ \xrightarrow{\quad} Z+1 \end{array}$$

$g \circ f$

$$X \begin{array}{c} \xrightarrow{f} Y \xrightarrow{k_1} Y+1 \\ \xrightarrow{\quad} Y+1 \end{array}$$

a partial map $X \rightarrow Y$

composition of partial maps

a map as a (total) partial map

Predicates and partial maps

$$p: X \rightarrow 1+1$$

$$X \xrightarrow{!} 1 \xrightarrow{k_1} 1+1$$

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predicate on X

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$$X \xrightarrow{f} Y+1 \xrightarrow{!+id} 1+1$$

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$\hat{!} \circ f$ is a measure of partiality of f

Predicates and partial maps

$$p: X \rightarrow 1+1$$

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predicate on X

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a partial map $X \rightarrow Y$

composition of partial maps

a map as a (total) partial map

$\hat{1} \circ f$ is a measure of partiality of f

$$\hat{1} \circ f = 1 \Rightarrow f \text{ total}$$

$$\hat{1} \circ f = 0 \Rightarrow f = \kappa_2!$$

Scalars and states

$$\lambda: 1 \rightarrow 1+1$$

a scalar (so partial map $1 \rightarrow 1$)

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$$\lambda \circ \mu \equiv \lambda \hat{\circ} \mu$$

product is composition as partial maps

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$\lambda \circ \mu \equiv \lambda \hat{\circ} \mu$ product is composition as partial maps

Similarly $\lambda \cdot p \equiv \lambda \hat{\circ} p$ for scalar $\lambda: 1 \rightarrow 1$

predicate $p: X \rightarrow 1$

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$\omega: 1 \rightarrow X$ a state on X

Scalars and states

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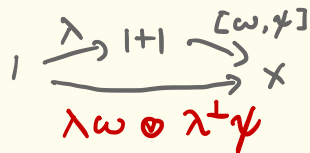
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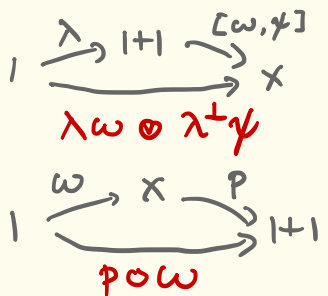
a state on X



convex combination of states

Scalars and states

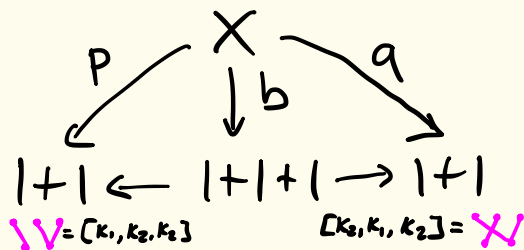
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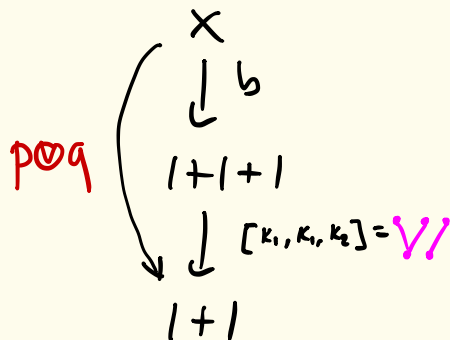
 $\lambda \omega \oplus \lambda^\perp \psi$
 $p \circ \omega$
 convex combination of states
 validity of a predicate in a state

Partial sum of predicates

If for predicates p, q on X , there is a b with



then p and q are **summable** (in symbols: $p \perp q$)
and their sum $p \oplus q$ is given by



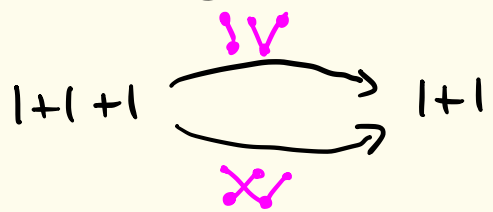
Definition an **Effectus** is a category \mathcal{C} with

- finite coproducts and final object 1 ,
- where all diagrams of the form

$$\begin{array}{ccc}
 x+y & \xrightarrow{id+!} & x+1 \\
 \downarrow !+id & & \downarrow !+id \\
 1+y & \xrightarrow{id+!} & 1+1
 \end{array}
 \qquad
 \begin{array}{ccc}
 x & \xrightarrow{!} & 1 \\
 \downarrow k_1 & & \downarrow k_1 \\
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 \end{array}$$

are pullback squares and

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 1 + 1 + 1 & \begin{array}{c} \xrightarrow{!+!} \\ \xrightarrow{!+!} \end{array} & 1 + 1
 \end{array}$$

Examples: \mathbf{vN}^{op} , \mathbf{Set} , $\mathbf{CRng}^{\text{op}}$, $\mathbf{Kl}(D)$, any topos, \mathbf{EA}^{op} , ...

Structure in an effectus

M , the set of scalars is an **effect monoid**,
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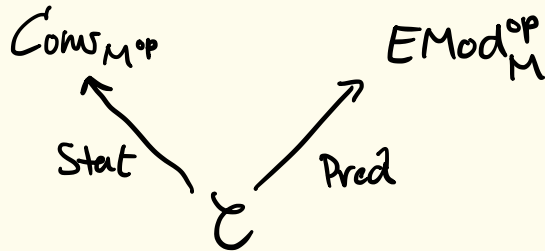
$\text{Stat } X$, the set of states on X is an (abstract) **M^{op} -convex set**.

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Definition a category C is an **Effectus** in partial form if

1. C is a fin PAC — that is

a. C has coproducts

b. C is PCM-enriched, i.e.:

α. every $\text{Hom}(X, Y)$ has partial binary operation \circledast and distinguished map 0 that turn it into a PCM

β. $f \perp g \Rightarrow \begin{cases} h \circ f \perp h \circ g & (h \circ f) \circledast (h \circ g) = h \circ (f \circledast g) \\ f \circ k \perp g \circ k & (f \circ k) \circledast (g \circ k) = (f \circledast g) \circ k \end{cases}$

c. $\triangleright_1 \circ b \perp \triangleright_2 \circ b$ for any $b: X \rightarrow Y + Y$,

where $\triangleright_1 \equiv [id, 0]: Y + Y \rightarrow Y$ and $\triangleright_2 \equiv [0, id]$.

d. $f \perp g \Rightarrow k_1 \circ f \perp k_2 \circ g$

2. C "has effects" — that is: there is an object I such that

a. the PCM $\text{Hom}(X, I) \equiv \text{Pred } X$ is an effect algebra $\forall X$.

b. $1 \circ f \perp 1 \circ g \Rightarrow f \perp g$

c. $1 \circ f = 0 \Rightarrow f = 0$.

A map f is called **total** iff $1 \circ f = 1$

Definition a category C is an **Effectus** in partial form if

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b. C is PCM —

α. every H_c
and d^*

β. $f \perp g \Rightarrow$

c. $\triangleright_1 \circ b \perp \triangleright_2 \circ$

where $\triangleright_1 \equiv$

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$(M, \oplus, 0)$ with $\oplus: M^2 \rightarrow M$,
 $0 \in M$ is a PCM iff

- $x \perp y \Rightarrow \begin{cases} y \perp x \\ x \oplus y = y \oplus x \end{cases}$
- $0 \perp x$ and $0 \oplus x = x$
- $x \perp y$ and $x \oplus y \perp z$
 $\Rightarrow \begin{cases} y \perp z, x \perp y \oplus z \\ (x \oplus y) \oplus z = x \oplus (y \oplus z) \end{cases}$

peration \oplus
into a PCM
 $\perp = k_0(f \oplus g)$
 $= (f \oplus g) \circ k$

$\triangleright_2 \equiv [0, id]$.

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c. $D_1 \circ b \perp D_2 \circ$

where $D_i \equiv$

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peration \oplus
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 $= (f \oplus g) \circ k$

$\dashv \circ \equiv [0, id]$.

2. C

a.

b.

c.

A PCM is an **Effect Algebra** iff

- for every x , there is a unique x^\perp such that

$$x^\perp \oplus x = 1 \equiv 0^\perp$$

- $x \perp 1 \Rightarrow x = 0$

object I such that
an effect algebra $\forall X$.

A map f

They're the "same"

If C is an effectus, then $\text{Par } C$, its category of partial maps, is an effectus in partial form.

If D is an effectus in partial form, then $\text{Tot } D$, its subcategory of total maps is an effectus.

Furthermore $\text{Par Tot } D \cong D$
and $\text{Tot Par } C \cong C$.

◇-effectus

▷-effectus preparation: quotients

An effectus (in partial form) has quotients if:

For every predicate $p: X \rightarrow 1$, there is an obj. X/p^\perp and (partial) map $\xi_{p^\perp}: X \rightarrow X/p^\perp$ with $1 \circ \xi_{p^\perp} \leq p$, such that for any other $f: X \rightarrow Y$ with $1 \circ f \leq p$, there is a unique g with

$$\begin{array}{ccc} X & \xrightarrow{\xi_{p^\perp}} & X/p^\perp \\ & \searrow f & \vdash g \\ & & Y \end{array}$$

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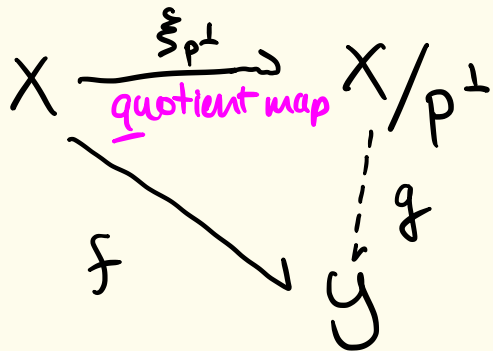
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(In vN , $\xi: \uparrow_p \downarrow A \uparrow_p \rightarrow A$
given by $\xi(a) = \sqrt{p} a \sqrt{p}$)

◇-effectus preparation: comprehension

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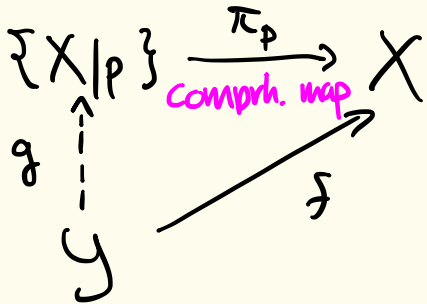
For every predicate $p: X \rightarrow 1$, there is an obj. $\{X|p\}$ and (partial) map $\pi_p: \{X|p\} \rightarrow X$ with $p \circ \pi_p = 1 \circ \pi_p$ such that for any other $f: Y \rightarrow X$ with $p \circ f = 1 \circ f$, there is a unique g with

$$\begin{array}{ccc} \{X|p\} & \xrightarrow{\pi_p} & X \\ \uparrow g & \nearrow f & \\ Y & & \end{array}$$

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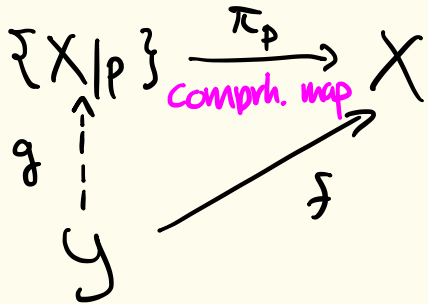
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(In vN , $\pi: A \rightarrow \lfloor p \rfloor A \lfloor p \rfloor$
given by $\pi(a) = \lfloor p \rfloor a \lfloor p \rfloor$)

Aside: why these names?

For effectus C , define category $\int \text{Pred}_{\square}$ by

- objects: pairs (X, p) , X object in C , $p \in \text{Pred } X$
- an arrow $(X, p) \rightarrow (Y, q)$
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There is an obvious $u: \int \text{Pred}_{\square} \rightarrow \text{Par } C$, $(X, p) \mapsto X$
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$$\begin{array}{ccc}
 & \int \text{Pred}_{\square} & \\
 0 \nearrow & \downarrow u & \searrow 1 \\
 & \text{Par } C &
 \end{array}$$

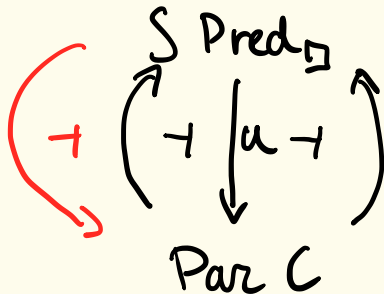
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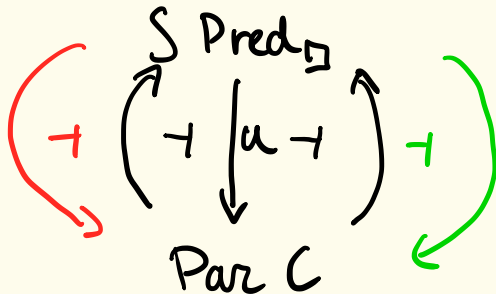
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▷-effectus preparation: images & sharp predicates

An effectus (in partial form) has images if:

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(In vN : sharp iff projection)

Dfn. a \diamond -effectus is an effectus
with quotients,
comprehension
and images
with: s sharp $\Rightarrow s^\perp$ sharp.

Proposition. In a \square -effectus the sharp predicates $\mathcal{S} \text{Pred } X$ on X are a sub-effect algebra of $\text{Pred } X$ and an orthomodular lattice.

Factorization in \mathbb{D} -effectus

$$X \xrightarrow{f} Y$$

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The pure maps $B(H) \rightarrow B(K)$ are

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Theorem. an ncp-map $\varphi: A \rightarrow B$
with Paschke / Stinespring dilation

$$\begin{array}{ccc} A & \xrightarrow{\varphi} & B \\ & \searrow \rho & \nearrow h \\ & \mathcal{K} & \end{array}$$

is pure if and only if ρ is surjective.

Ceiling $\lceil p \rceil$ and floor $\lfloor p \rfloor$

Define $\lfloor p \rfloor \equiv \text{im } \pi_p$ (π_p comprh. for p)

$$\lceil p \rceil \equiv \lfloor p \rfloor^\perp$$

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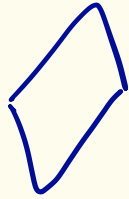
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(In vN : $\lceil f(a) \rceil = \lceil f(\lceil a \rceil) \rceil$ useful rule)



\square , the possibilistic restriction

For $f: X \rightarrow Y$ in a \square -effectus, define

$$\text{SPred } X \begin{array}{c} \xrightarrow{f^\square} \\ \xleftarrow{f^\square} \end{array} \text{SPred } Y$$

by

$$f^\square(s) \equiv \ulcorner s \circ f \urcorner$$

$$f^\square(s) \equiv \text{im}(f \circ \pi_s)$$

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In vN :

$$f^\square = g^\square \text{ iff}$$

for every normal
state ω and effect a ,
we have

$$\omega(f(a)) = 0 \Leftrightarrow \omega(g(a))$$

Proposition. In a \triangleleft -effectus:

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$$\bullet \text{im } f = f_\circ(1) \quad \ulcorner 1 \circ f \urcorner = f^\circ(1)$$

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$\left[\begin{array}{l} f \text{ is pure} \\ f = g \circ g \text{ for some } \diamond\text{-self adjoint } g. \end{array} \right.$

Theorem. In vN , the \diamond -positive maps are precisely $a \mapsto \sqrt{b} a \sqrt{b}$.

Δ -effectus

Dfn. an \diamond -effectus is \diamond -effectus where

- for every predicate p on X , there is a unique \diamond -positive map $\text{asrt}_p : X \rightarrow X$ with $1 \circ \text{asrt}_p = p$.
- $\xi \circ \pi$ is pure for any comprehension π and quotient map ξ .

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Write $p \Delta q \equiv q \circ \text{asrt}_p$ and $p^2 \equiv p \Delta p$.

Proposition. In an $\&$ -effectus TFAE

- p is sharp
- $p \& p = p$
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Proposition. In an $\&$ -effectus

$$\text{im } f \leq s \quad \Leftrightarrow \quad \text{asrt}_s \circ f = f$$

$$1 \circ f \leq t \quad \Leftrightarrow \quad f \circ \text{asrt}_t = f$$

Polar decomposition

In an \diamond -effectus, any pure map $f: X \rightarrow Y$ factors as follows:

$$X \xrightarrow{\text{asrt}_{1of}} X \xrightarrow{\xi} X_{/1of} \xrightarrow{\approx} \{y | \text{im } f\} \xrightarrow{\pi} Y$$

\uparrow
 \diamond -positive

pure map h with $1oh$ sharp
 \approx partial isometry

+ effectus

Dfn. a \dagger -effectus \mathcal{C} is an $\&$ -effectus, where the subcat. of pure maps is a \dagger -category with:

- $\text{asrt}_p^\dagger = \text{asrt}_p$
- f is \diamond -adjoint to f^\dagger
- for every \dagger -positive f , there is a unique \dagger -positive g with $f = g \circ g$.
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It follows $\lceil 1 \circ f^\dagger \rceil = \text{im} f$ and $\text{im} f^\dagger = \lceil 1 \circ f \rceil$,

(so with slight abuse of notation:)

$$\pi_S^\dagger = \sum_{S^\perp} \quad \text{and} \quad \sum_{S^\perp}^\dagger = \pi_S \quad (\text{sharp } S)$$

Theorem. An $\&$ -effectus is a \dagger -effectus iff

- for every predicate p ,
there is a unique predicate q
with $q \& q = p$
- $\text{asrt}_{p \& q}^2 = \text{asrt}_p \circ \text{asrt}_q^2 \circ \text{asrt}_p$
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↪ Cf. Fundamental Id. of Quad. Jordan Algs.

- $t \circ \xi_s$ is sharp for all sharp s, t .

Dfn. An effectus is **operational** iff

- the scalars are isomorphic to $[0, 1]$,
- the predicates are jointly monic,
- $p \leq q \iff \forall \text{ state } \omega. p \circ \omega \leq q \circ \omega$,
- every object X is 'finite-dimensional':
that is: $\text{Stat } X$ 'is' a closed convex subset of a finite-dimensional vector space.

Theorem. (Wettering)

- The category EJA of Euclidean Jordan Algebras with positive maps is a \dagger -effectus.
- Any operational \dagger -effectus is equivalent to a subcategory of EJA .

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(vN^{op} is \dagger -effectus, but not operational)

Theorem. Every t -effectus is a
homological category in the sense
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Corollary. Grandis' Snake Lemma
holds for von Neumann algebras...

Grandis' Snake Lemma. If we have a diagram

$$\begin{array}{ccccccc}
 A & \xrightarrow{f} & B & \xrightarrow{g} & C & \longrightarrow & 0 \\
 & & \downarrow a & & \downarrow b & & \downarrow c \\
 0 & \longrightarrow & A' & \xrightarrow{h} & B' & \xrightarrow{k} & C'
 \end{array}$$

in a \dagger -effectus such that

- $\text{im } f = \Gamma_{1og}^\perp$
- $\text{im } h = \Gamma_{1ok}^\perp$
- g quotient with $1og$ sharp
- h comprehension

- $b^\square(b_\diamond(\text{im } f)) = \Gamma_{1ob}^\perp \vee \text{im } f$
- $b_\diamond(b^\square(\text{im } h)) = (\text{im } h) \wedge \text{im } b$
- $k^\square(k_\diamond(\text{im } b)) = (\text{im } h) \vee \text{im } b$
- $f_\diamond(f^\square(b^\square(0))) = \Gamma_{1ob}^\perp \wedge \text{im } f$

then ...

$$\{A | (10a)^\perp\} \xrightarrow{\bar{f}} \{B | (10b)^\perp\} \xrightarrow{\bar{g}} \{C | (10c)^\perp\} \quad \exists d$$

$$\begin{array}{ccc} \downarrow \pi_{(10a)^\perp} & & \downarrow \pi_{(10b)^\perp} & & \downarrow \pi_{(10c)^\perp} \\ A & \xrightarrow{f} & B & \xrightarrow{g} & C & \longrightarrow & 0 \end{array}$$

$$\begin{array}{ccc} \downarrow a & \xrightarrow{\quad} & \downarrow b & \xrightarrow{\quad} & \downarrow c \end{array}$$

$$0 \longrightarrow A' \xrightarrow{h} B' \xrightarrow{k} C'$$

$$\begin{array}{ccc} \downarrow \xi_{\text{ima}} & & \downarrow \xi_{\text{imb}} & & \downarrow \xi_{\text{imc}} \end{array}$$

$$\begin{array}{ccc} \longrightarrow A'/\text{ima} & \xrightarrow{h} & B'/\text{imb} & \xrightarrow{k} & C'/\text{imc} \end{array}$$

with $\text{im } \bar{f} = \pi_{\text{ro } \bar{g}}^\perp$, $\text{im } \bar{g} = \pi_{\text{ro } \bar{h}}^\perp$, ...

Take away

- QT reconstructions don't need dilations/purifications or parallel composition \otimes .
- $\text{comprh} \circ \text{iso} \circ \text{quot.}$ is the right kind of pure.
- $()^\diamond$, \Downarrow -adjoint and \Downarrow -positive.

Further reading

- "Dagger and dilations", BW [arXiv 1803.01911](https://arxiv.org/abs/1803.01911).
- "Introduction to Effectus theory"
Cho, Jacobs, Westerbaan, BW [arXiv 1512.05813](https://arxiv.org/abs/1512.05813)
- "Reconstruction of Quantum Theory from univ. filters"
Wetering [arXiv 1801.05796](https://arxiv.org/abs/1801.05796)