

HOMOTOPY-COHERENT ALGEBRAS AND POLYNOMIAL MONADS

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Let \mathbb{F}_* = category of pointed finite sets and let \mathcal{S} = ∞ -category of spaces. A special Γ -space is a functor

$$F: \mathbb{F}_* \rightarrow \mathcal{S}$$

satisfying the *Segal condition* ($F(\langle n \rangle) \xrightarrow{\sim} \prod_{i=1}^n F(\langle 1 \rangle)$).

Examples

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- ∞ -operads: the dendroidal category Ω^{op} of Moerdijk and Weiss
- ∞ -properads: Γ^{op} of Hackney, Robertson and Yau

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But in which generality can we talk about “Segal conditions”?

We observe that all the examples have certain key features in common. This leads us to the following definition.

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An **algebraic pattern** \mathcal{P} is an ∞ -category equipped with

- an inert-active factorization system $(\mathcal{P}^{\text{int}}, \mathcal{P}^{\text{act}})$,
- a full subcategory $\mathcal{P}^{\text{el}} \subseteq \mathcal{P}^{\text{int}}$ (its objects are called *elementary*).

Definition

Let \mathcal{P} be an algebraic pattern and let $\mathcal{P}_{X/}^{\text{el}} := \mathcal{P}^{\text{el}} \times_{\mathcal{P}^{\text{int}}} \mathcal{P}_{X/}^{\text{int}}$.

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Then a *Segal \mathcal{P} -space* or *object* is a functor $F: \mathcal{P} \rightarrow \mathcal{S}$ (or $\mathcal{C} = \text{Cat}_{\infty}, \infty\text{-topos}, \dots$) which satisfies the “Segal condition”, i.e.

$$F(X) \xrightarrow{\sim} \lim_{E \in \mathcal{P}_{X/}^{\text{el}}} F(E).$$

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- elementary objects \xrightarrow{F} building blocks
- inert maps \xrightarrow{F} geometric operations = “restriction maps” which are used for the Segal condition
- active maps \xrightarrow{F} algebraic operations (multiplications, compositions,...)

Notation

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- Ξ = category of unrooted trees defined by HRY. \Rightarrow Cyclic ∞ -operads = Segal Ξ^{op} -spaces.

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- Ξ = category of unrooted trees defined by HRY. \Rightarrow Cyclic ∞ -operads = Segal Ξ^{op} -spaces.
- \mathcal{O} = Lurie's (generalized) ∞ -operad $\Rightarrow \text{Alg}_{\mathcal{O}}(\mathcal{C}) = \text{Segal } \mathcal{O}\text{-objects in } \mathcal{C}$.

Remark

A functor $f: \mathcal{P} \rightarrow \mathcal{Q}$ induces adjunctions

$$f_! : \text{Fun}(\mathcal{P}, \mathcal{S}) \rightleftarrows \text{Fun}(\mathcal{Q}, \mathcal{S}) : f^* \text{ and } f^* : \text{Fun}(\mathcal{Q}, \mathcal{S}) \rightleftarrows \text{Fun}(\mathcal{P}, \mathcal{S}) : f_*.$$

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Under certain *checkable* conditions f^* , $f_!$ and f_* restrict to functors

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\mathcal{P} is *extendable* if $i_! : \text{Fun}(\mathcal{P}^{\text{el}}, \mathcal{S}) \simeq \text{Seg}_{\mathcal{P}^{\text{int}}}(\mathcal{S}) \rightarrow \text{Seg}_{\mathcal{P}}(\mathcal{S})$ is given by restriction, where $i : \mathcal{P}^{\text{int}} \hookrightarrow \mathcal{P}$ is the inclusion.

Applications

By checking conditions

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By checking conditions

- for right Kan extensions f_* : The forgetful functor $\text{CycOpd}_\infty \rightarrow \text{Opd}_\infty$ has a right adjoint.
- for left Kan extensions $f_!$: We recover the formula for operadic left Kan extensions in the sense of Lurie.
- for extendability: Θ_n^{op} , $\Delta^{n,\text{op}}$ and Ω^{op} are extendable. \Rightarrow Formula for free (∞, n) -categories, free n -fold ∞ -categories and free ∞ -operads.

Proposition

For every \mathcal{P} the adjunction $\text{Fun}(\mathcal{P}^{\text{el}}, \mathcal{S}) \rightleftarrows \text{Seg}_{\mathcal{P}}(\mathcal{S})$ is monadic.

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Questions

Which monads on presheaf ∞ -categories can be described as the free Segal \mathcal{P} -space monad for an extendable pattern \mathcal{P} ? How far is this correspondence from being an equivalence?

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 - where an object is a polynomial monadic right adjoints over some functor category $\text{Fun}(\mathcal{I}, \mathcal{S})$, (a monad is *polynomial* if it is cartesian and a local right adjoint)
 - where a morphism is a commutative square whose mate transformation is cartesian.

Theorem

If \mathcal{P} is extendable, then the associated monad $T_{\mathcal{P}}$ on $\mathbf{Fun}(\mathcal{P}^{\text{el}}, \mathcal{S})$ is polynomial. Hence, we have a functor

$$\mathfrak{M}: \text{ExtPatt} \rightarrow \text{PolyMnd}.$$

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We now want to prove that this functor is essentially surjective. For a given polynomial monad we want to construct the associated algebraic pattern.

Remark

The main input for the proof for the essential surjectivity is an ∞ -categorical version of work of Berger, Melliés and Weber and it is closely related to the “nerve theorem” studied by Leinster, Kock and Weber.

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The difference between cartesian monads (induced by Σ -free operads) and *weakly* cartesian monads in the 1-categorical world vanishes by going to that of ∞ -categories.

Construction

Given a polynomial monad T on the presheaf ∞ -category $\mathbf{Fun}(\mathcal{I}, \mathcal{S})$, define \mathcal{P}_T by

$$\begin{array}{ccccc}
 \mathcal{P}_T^{\text{op}} & \hookrightarrow & \mathbf{Alg}_T(\mathbf{Fun}(\mathcal{I}, \mathcal{S})) & & \\
 \uparrow F_T & & & & \uparrow F_T \downarrow U_T \\
 \mathcal{I}^{\text{op}} & \hookrightarrow & \mathcal{P}_T^{\text{int,op}} & \hookrightarrow & \mathbf{Fun}(\mathcal{I}, \mathcal{S})
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The construction defines an algebraic pattern \mathcal{P}_T with a factorization system $(\mathcal{P}_T^{\text{int}}, \mathcal{P}_T^{\text{act}})$ and $\mathcal{P}_T^{\text{el}} = \mathcal{I}$.

Theorem

The assignment $T \mapsto \mathcal{P}_T$ gives a functor

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Moreover, this functor is fully faithful and $\mathfrak{M}\mathfrak{P} \simeq \text{id}$.

Examples

- Free operad monad on $\mathbf{Fun}(\mathcal{I}, \mathcal{S})$ (\mathcal{I} = the category of trees with at most one vertex) $\overset{\mathfrak{F}}{\mapsto} \Omega^{\text{op}}$.

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What is the essential image of the fully faithful functor $\mathfrak{P}: \text{PolyMnd} \rightarrow \text{ExtPatt}$?

Theorem

An extendable algebraic pattern lies in the essential image of \mathfrak{P} iff it is *nice*, i.e. every object $X \in \mathcal{P}$ admits an active map $X \rightarrow E$, $E \in \mathcal{P}^{\text{el}}$ and $\text{Map}_{\mathcal{P}^{\text{int}}}(X, -) \in \text{Seg}_{\mathcal{P}^{\text{int}}}(\mathcal{S})$.

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Many algebraic patterns are nice: $\Delta^{n, \text{op}}$, Θ_n^{op} , Ω^{op} , \dots

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\mathbb{F}_* is not nice, $\mathfrak{P}\mathfrak{M}(\mathbb{F}_*) \simeq \text{Span}^{\text{inj}}(\mathbb{F})$.

ExtPatt_n = the full subcategory of ExtPatt spanned by nice algebraic patterns.

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Corollary

The adjunction $\mathfrak{M} : \text{ExtPatt} \rightleftarrows \text{PolyMnd} : \mathfrak{P}$ restricts to an equivalence

$$\text{ExtPatt}_n \simeq \text{PolyMnd}.$$

In particular, PolyMnd is a localization of ExtPatt .