

# Retrocells

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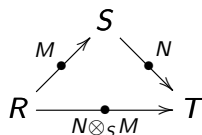
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# Bimodules

- The bicategory  $\mathcal{Bim}$  has rings  $R, S, T, \dots$  as objects, bimodules  $M : R \rightarrow S$  as 1-cells, and  $S$ - $R$ -linear maps as 2-cells

Composition is  $\otimes$



- $\mathcal{Bim}$  is biclosed,  $\otimes$  has right adjoints in each variable

$$\underline{M \rightarrow N \otimes_T P}$$

$$\underline{N \otimes_S M \rightarrow P}$$

$$N \rightarrow P \otimes_R M$$

$$N \otimes_T P = \text{Hom}_T(N, P), \quad P \otimes_R M = \text{Hom}_R(M, P)$$

Many bicategories are biclosed

- $Bim$  : Rings, bimodules, linear maps
- $Prof$  : Categories, profunctors, natural transformations
- $\mathbf{V}\text{-}Prof$  :  $\mathbf{V}$  – with colimits preserved by  $\otimes$ 
  - biclosed
  - limits
- $Span(\mathbf{A})$  :  $\mathbf{A}$  with pullbacks and locally cartesian closed

# Scandal

Good bicategories (all of the above) are the vertical part of naturally occurring double categories:

Ring, Cat, **V-Cat**, Span**A**

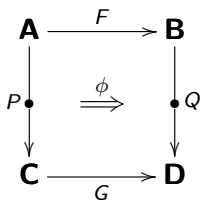
But the internal homs  $\otimes$  and  $\boxtimes$  are not double functors!

# Double categories

- A *double category* is a “category with two sorts of morphisms”
- **Example:**  $\mathbb{R}\text{ing}$

$$\begin{array}{ccc} R & \xrightarrow{f} & S \\ \downarrow M & \xRightarrow{\alpha} & \downarrow N \\ R' & \xrightarrow{f'} & S' \end{array}$$

- Example:  $\mathbf{Cat}$

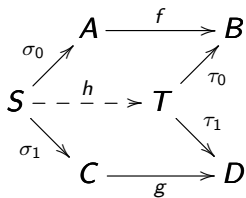


$$P : \mathbf{A}^{op} \times \mathbf{C} \longrightarrow \mathbf{Set}$$

$$Q : \mathbf{B}^{op} \times \mathbf{D} \longrightarrow \mathbf{Set}$$

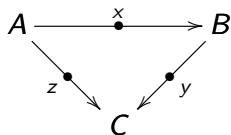
$$\phi : P(-, =) \longrightarrow Q(F-, G=)$$

- **Example:** Span **A**



## Left homs

- $\mathbb{A}$  has *left homs* if  $y \bullet ( )$  has a right adjoint  $y \blacktriangleright ( )$  in  $\mathcal{V}er\mathbb{A}$



$$\frac{y \bullet x \rightarrow z}{x \rightarrow y \blacktriangleright z}$$

in  $\mathcal{V}er\mathbb{A}$

Mike Shulman, “Framed bicategories and monoidal fibrations” (TAC 2008)  
Roald Koudenburg, “On pointwise Kan extensions in double categories”  
(TAC 2014)



## Respecting boundaries

- $y \backslash z$  is covariant in  $z$  and contravariant in  $y$

$$y' \xrightarrow{\beta} y, z \xrightarrow{\gamma} z' \quad \rightsquigarrow \quad y \backslash z \xrightarrow{\beta \backslash \gamma} y' \backslash z'$$

## Respecting boundaries

- $y \multimap z$  is covariant in  $z$  and contravariant in  $y$

$$y' \xrightarrow{\beta} y, z \xrightarrow{\gamma} z' \rightsquigarrow y \multimap z \xrightarrow{\beta \multimap \gamma} y' \multimap z'$$

- We have evaluation  $\epsilon : y \bullet (y \multimap z) \rightarrow y$

$$\begin{array}{c}
 \begin{array}{cccc}
 A & = & A & = & A & = & A \\
 \downarrow y \multimap z & & \downarrow y \multimap z & & \downarrow & & \downarrow \\
 B & = & B & \xRightarrow{\epsilon} & z & \xRightarrow{\gamma} & z' \\
 \downarrow y' & \xRightarrow{\beta} & \downarrow y & & \downarrow & & \downarrow \\
 C & = & C & = & C & = & C
 \end{array} \\
 \hline
 y \multimap z \rightarrow y' \multimap z'
 \end{array}$$



# Globular universal

$$\forall \quad \begin{array}{ccc} A & \equiv & A \\ \downarrow x & & \downarrow \\ \bullet & & \\ \downarrow & \Rightarrow^Q & \downarrow z \\ B & & \bullet \\ \downarrow y & & \downarrow \\ \bullet & & \\ \downarrow & & \downarrow \\ C & \equiv & C \end{array} \quad \exists! \quad \begin{array}{ccc} A & \equiv & A \\ \downarrow x & \Rightarrow^\beta & \downarrow y \quad z \\ \bullet & & \bullet \\ \downarrow & & \downarrow \\ B & \equiv & B \end{array}$$

s.t.

$$\begin{array}{ccc} A & \equiv & A & \equiv & A \\ \downarrow x & \Rightarrow^\beta & \downarrow y \quad z & & \downarrow \\ \bullet & & \bullet & & \\ \downarrow & & \downarrow & \Rightarrow^\epsilon & \downarrow z \\ B & \equiv & B & & \bullet \\ \downarrow y & = & \downarrow y & & \downarrow \\ \bullet & & \bullet & & \\ \downarrow & & \downarrow & & \downarrow \\ C & \equiv & C & \equiv & C \end{array} = \begin{array}{ccc} A & \equiv & A \\ \downarrow x & & \downarrow \\ \bullet & & \\ \downarrow & \Rightarrow^Q & \downarrow z \\ B & & \bullet \\ \downarrow y & & \downarrow \\ \bullet & & \\ \downarrow & & \downarrow \\ C & \equiv & C \end{array}$$

# More universal

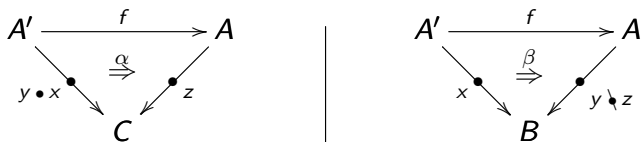
$$\forall \begin{array}{ccc} A' & \xrightarrow{f} & A \\ \downarrow x & & \downarrow \\ B & \xrightarrow{\alpha} & z \\ \downarrow y & & \downarrow \\ C & \xlongequal{\quad} & C \end{array} \qquad \exists! \begin{array}{ccc} A' & \xrightarrow{f} & A \\ \downarrow x & \xRightarrow{\beta} & \downarrow y \quad z \\ B & \xlongequal{\quad} & B \end{array}$$

s.t.

$$\begin{array}{ccc} A' & \xrightarrow{f} & A \xlongequal{\quad} A \\ \downarrow x & \xRightarrow{\beta} & \downarrow y \quad z \\ B & \xlongequal{\quad} B & \xrightarrow{\epsilon} z \\ \downarrow y & \xlongequal{\quad} \downarrow y & \downarrow \\ C & \xlongequal{\quad} C \xlongequal{\quad} C \end{array} = \begin{array}{ccc} A' & \xrightarrow{f} & A \\ \downarrow x & & \downarrow \\ B & \xrightarrow{\alpha} & z \\ \downarrow y & & \downarrow \\ C & \xlongequal{\quad} & C \end{array}$$

# Strong universality

*Strong universal property:*



# Companions

- In a double category  $\mathbb{A}$ , a vertical arrow  $v : A \bullet \rightarrow B$  is a *companion* of a horizontal arrow  $f : A \rightarrow B$  if there are *binding cells*  $\alpha$  and  $\beta$  such that

$$\begin{array}{ccc}
 A & \xrightarrow{1_A} & A & \xrightarrow{f} & B \\
 \text{id}_A \downarrow & \xRightarrow{\alpha} & \downarrow v & \xRightarrow{\beta} & \downarrow \text{id}_B \\
 A & \xrightarrow{f} & B & \xrightarrow{1_B} & B
 \end{array}
 =
 \begin{array}{ccc}
 A & \xrightarrow{f} & B \\
 \text{id}_A \downarrow & \xRightarrow{\text{id}_f} & \downarrow \text{id}_B \\
 A & \xrightarrow{f} & B
 \end{array}
 \quad \beta\alpha = \text{id}_f$$

$$\begin{array}{ccc}
 A & \xrightarrow{1_A} & A \\
 \text{id}_A \downarrow & \xRightarrow{\alpha} & \downarrow v \\
 A & \xrightarrow{f} & B \\
 v \downarrow & \xRightarrow{\beta} & \downarrow \text{id}_B \\
 B & \xrightarrow{1_B} & B
 \end{array}
 =
 \begin{array}{ccc}
 A & \xrightarrow{1_A} & A \\
 v \downarrow & \xRightarrow{1_v} & \downarrow v \\
 B & \xrightarrow{1_B} & B
 \end{array}
 \quad \beta \bullet \alpha = 1_v$$

# Properties

- Companions, when they exist, are unique up to globular isomorphism
- We make a choice of companion  $f_*$  and, following Ronnie Brown, denote the binding cells by corner brackets
- We have  $(1_A)_* \cong \text{id}_A$  and  $(gf)_* \cong g_*f_*$
  -

$$\begin{array}{ccc}
 \begin{array}{ccc}
 A & \xrightarrow{f} & B \\
 \downarrow v & \Downarrow \phi & \downarrow w \\
 C & \xrightarrow{g} & D
 \end{array} & \longmapsto & \begin{array}{ccc}
 A & \xlongequal{\quad} & A \\
 \parallel & \lrcorner & \downarrow f_* \\
 A & \xrightarrow{f} & B \\
 \downarrow v & \Downarrow \phi & \downarrow w \\
 C & \xrightarrow{g} & D \\
 \downarrow g_* & \llcorner & \parallel \\
 D & \xlongequal{\quad} & D
 \end{array} = \begin{array}{ccc}
 A & \xlongequal{\quad} & A \\
 \downarrow v & & \downarrow f_* \\
 C & \xRightarrow{\psi} & B \\
 \downarrow g_* & & \downarrow w \\
 D & \xlongequal{\quad} & D
 \end{array}
 \end{array}$$

gives a bijection between  $\phi$ 's and  $\psi$ 's



# Conjoints

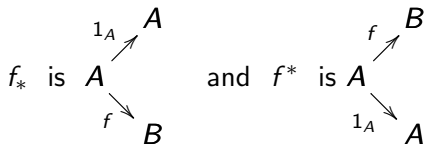
There is a dual notion of *conjoint*  $f^*$

$$\begin{array}{ccc}
 A & \xrightarrow{f} & B \xrightarrow{1_B} B \\
 \text{id}_A \downarrow & \xRightarrow{\psi} & \downarrow f^* \xRightarrow{\chi} \downarrow \text{id}_B \\
 A & \xrightarrow{1_A} & A \xrightarrow{f} B
 \end{array}
 =
 \begin{array}{ccc}
 A & \xrightarrow{f} & B \\
 \text{id}_A \downarrow & \xRightarrow{\text{id}_f} & \downarrow \text{id}_B \\
 A & \xrightarrow{f} & B
 \end{array}
 \quad \chi\psi = \text{id}_f$$

$$\begin{array}{ccc}
 B & \xrightarrow{1_B} & B \\
 f^* \downarrow & \xRightarrow{\chi} & \downarrow \text{id}_B \\
 A & \xrightarrow{f} & B \\
 \text{id}_A \downarrow & \xRightarrow{\psi} & \downarrow f^* \\
 A & \xrightarrow{1_A} & A
 \end{array}
 =
 \begin{array}{ccc}
 B & \xrightarrow{1_B} & B \\
 f^* \downarrow & \xRightarrow{1_{f^*}} & \downarrow f^* \\
 A & \xrightarrow{1_A} & A
 \end{array}
 \quad \psi \bullet \chi = 1_{f^*}$$

# Examples

- In  $\mathbb{R}\text{ing}$ ,  $f : R \longrightarrow S$   
 $f_*$  is  $S$  considered as an  $S$ - $R$  bimodule  
 $f^*$  is  $S$  considered as an  $R$ - $S$  bimodule
- In  $\mathbb{C}\text{at}$ ,  $F : \mathbf{A} \longrightarrow \mathbf{B}$   
 $F_* = \mathbf{B}(F-, =)$  and  $F^* = \mathbf{B}(-, F=)$
- In  $\text{Span}(\mathbf{A})$ ,  $f : A \longrightarrow B$



## What strong means

- The strong universal property is equivalent to the globular one plus the stability property

$$y \multimap (z \bullet f_*) \cong (y \multimap z) \bullet f_*$$

- If every horizontal arrow has a conjoint, then the strong universal property is equivalent to the globular one

## Left duals

- Suppose  $\mathbb{A}$  left closed
- For  $v : A \multimap B$  we can define its *left dual*  $\bullet v = v \backslash \text{id}_B : B \multimap A$

We have

$$\bullet \text{id}_B \cong \text{id}_B$$

$$\bullet v \bullet w \longrightarrow \bullet(w \bullet v)$$

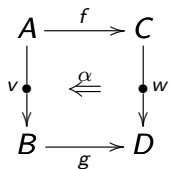
So perhaps we get a lax normal

$$\mathbb{A}^{\text{co}} \longrightarrow \mathbb{A}$$

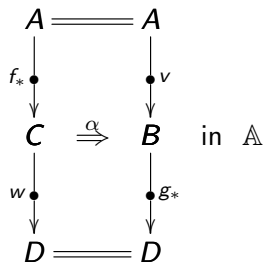
$$\begin{array}{ccc}
 A & \xrightarrow{f} & C \\
 \downarrow v & \xRightarrow{\alpha} & \downarrow w \\
 B & \xrightarrow{g} & D
 \end{array}
 \quad \rightsquigarrow \quad
 \begin{array}{ccc}
 B & \xrightarrow{g} & D \\
 \downarrow \bullet v & \xRightarrow{\bullet \alpha} & \downarrow \bullet w \\
 A & \xrightarrow{f} & C
 \end{array}$$

# Retrocells

A *retrocell*



is a cell



# Quintets

- **Example:** In  $\mathbb{Q}(\mathcal{A})$ , a cell is a quintet

$$\begin{array}{ccc} A & \xrightarrow{f} & B \\ h \downarrow & \swarrow & \downarrow k \\ C & \xrightarrow{g} & D \end{array}$$

and a retrocell is a coquintet

$$\begin{array}{ccc} A & \xrightarrow{f} & B \\ h \downarrow & \nearrow & \downarrow k \\ C & \xrightarrow{g} & C \end{array}$$

# Mates

## Proposition

(1) If  $v$  and  $w$  as below have right adjoints  $v'$  and  $w'$  in  $\mathcal{V}er\mathbb{A}$ , then retrocells  $\alpha$  are in bijection with standard cells  $\beta$ :

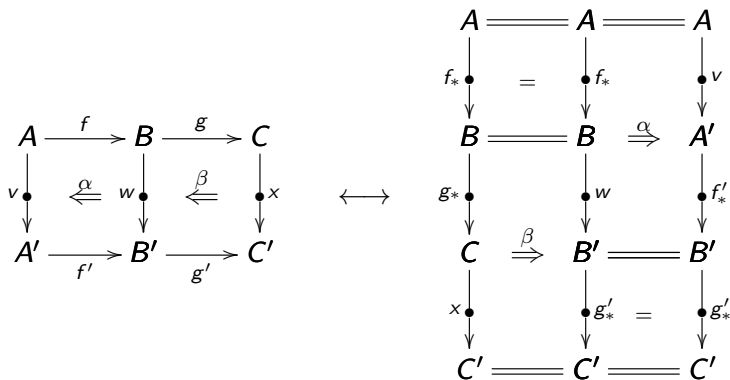
$$\begin{array}{ccc} A \xrightarrow{f} B & & C \xrightarrow{g} D \\ v \downarrow & \Leftarrow \alpha & \downarrow w \\ C \xrightarrow{g} D & & A \xrightarrow{f} B \end{array} \longleftrightarrow \begin{array}{ccc} C \xrightarrow{g} D & & A \xrightarrow{f} B \\ v' \downarrow & \Rightarrow \beta & \downarrow w' \\ A \xrightarrow{f} B & & C \xrightarrow{g} D \end{array}$$

(2) If  $f$  and  $g$  have right adjoints  $h$  and  $k$  in  $\mathcal{H}or\mathbb{A}$ , then retrocells  $\alpha$  are in bijection with standard cells  $\gamma$ :

$$\begin{array}{ccc} A \xrightarrow{f} B & & B \xrightarrow{h} A \\ v \downarrow & \Leftarrow \alpha & \downarrow w \\ C \xrightarrow{g} D & & D \xrightarrow{k} C \end{array} \longleftrightarrow \begin{array}{ccc} B \xrightarrow{h} A & & A \xrightarrow{f} B \\ w \downarrow & \Rightarrow \gamma & \downarrow v \\ D \xrightarrow{k} C & & C \xrightarrow{g} D \end{array}$$

# Composition

Retrocells can be composed horizontally





# Composition

and vertically

$$\begin{array}{ccc}
 A & \xrightarrow{f} & B \\
 \downarrow v & \xleftarrow{\alpha} & \downarrow w \\
 A' & \xrightarrow{f'} & B' \\
 \downarrow v' & \xleftarrow{\alpha'} & \downarrow w' \\
 A'' & \xrightarrow{f''} & B''
 \end{array}
 \quad \longleftrightarrow \quad
 \begin{array}{ccccc}
 A & \xlongequal{\quad} & A & \xlongequal{\quad} & A \\
 \downarrow f_* \bullet & & \downarrow v & = & \downarrow v \\
 B & \xrightarrow{\alpha} & A' & \xlongequal{\quad} & A' \\
 \downarrow w & & \downarrow f'_* \bullet & & \downarrow v' \\
 B' & \xlongequal{\quad} & B' & \xrightarrow{\alpha'} & A'' \\
 \downarrow w' \bullet & = & \downarrow w' & & \downarrow f''_* \bullet \\
 B'' & \xlongequal{\quad} & B'' & \xlongequal{\quad} & B''
 \end{array}$$

## Theorem

This gives a double category  $\mathbb{A}^{\text{ret}}$ .  $\mathbb{A}^{\text{ret}}$  has companions and  $(\mathbb{A}^{\text{ret}})^{\text{ret}} \cong \mathbb{A}$

## Commuter cells

- In M. Grandis, R. Paré, Kan extensions in double categories, TAC 2008, we introduced *commutative cells* to express the universal property of comma double categories

$$\begin{array}{ccc}
 A & \xrightarrow{f} & B \\
 \downarrow v & \Downarrow Q & \downarrow w \\
 C & \xrightarrow{g} & D
 \end{array}$$

is a *commuter cell* if

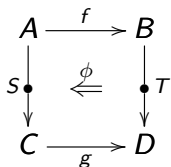
$$\begin{array}{ccc}
 A & \xlongequal{\quad} & A \\
 \parallel & \lrcorner & \downarrow f_* \\
 A & \xrightarrow{f} & B \\
 \downarrow v & \Downarrow Q & \downarrow w \\
 C & \xrightarrow{g} & D \\
 \downarrow g_* & \llcorner & \parallel \\
 D & \xlongequal{\quad} & D
 \end{array}$$

is horizontally invertible

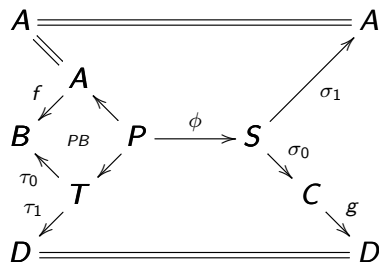
- The inverse would be a retrocell

# Retrocells of spans

- In  $\text{Span}(\mathbf{A})$



is



- In  $\text{Set} = \text{Span}(\mathbf{Set})$

Denote an element of  $S$  by  $s : a \bullet \rightarrow c$  ( $\sigma_0 s = a$ ,  $\sigma_1 s = c$ )

Then

$$\phi : (a, fa \xrightarrow{t} d) \mapsto (a \xrightarrow{\phi t} c_t), \quad g(c_t) = d$$

## Category objects

- A category object in  $\mathbf{A}$  is a vertical monad in  $\text{Span}(\mathbf{A})$
- An internal functor  $F : \mathbb{A} \rightarrow \mathbb{B}$  is a cell

$$\begin{array}{ccc}
 A_0 & \xrightarrow{F_0} & B_0 \\
 A_1 \downarrow \bullet & \xRightarrow{F_1} & \bullet \downarrow B_1 \\
 A_0 & \xrightarrow{F_0} & B_0
 \end{array}$$

respecting composition and identities

- A retrocell  $\phi$  is an object function  $F_0$  together with a lifting operation

$$\begin{array}{ccc}
 \mathbb{A} & & \\
 \vdots & & \\
 \mathbb{B} & & \\
 & \begin{array}{ccc}
 A & \overset{\phi b}{\dashrightarrow} & A_b \\
 \downarrow & & \downarrow \\
 F_0 A & \xrightarrow{b} & B
 \end{array} & & 
 \end{array}$$

- If  $\phi$  respects composition and identities, then this is exactly a *cofunctor*  $\mathbb{B} \rightarrow \mathbb{A}$  in the sense of Aguiar

## Discrete opfibrations

- $\phi$  looks like the lifting property for opfibrations without the projection functor
- If  $F$  is also a functor and  $F_1$  and  $\phi$  are companions in a certain double category of cells and retrocells, then  $F$  is a discrete opfibration. In fact  $F$  is a discrete opfibration if and only if  $F_1$  is a commuter cell

# Lax functors

- If  $F : \mathbb{A} \longrightarrow \mathbb{B}$  is a double functor, we get  $F^{ret} : \mathbb{A}^{ret} \longrightarrow \mathbb{B}^{ret}$
- If  $F : \mathbb{A} \longrightarrow \mathbb{B}$  is just lax, it doesn't extend to  $\mathbb{A}^{ret}$ ; it should properly respect companions
- If  $F$  is lax normal, then  $F$  preserves companions and also composites of the form  $A \xrightarrow{\bullet \text{ } f_*} B \xrightarrow{\bullet \text{ } v} C$

$$\phi(v, f_*) : F(v) \bullet F(f_*) \longrightarrow F(v \bullet f_*) \quad \text{iso}$$

[Dawson, Paré, Pronk, The Span Construction, TAC 2010]

# Paranormal

## Definition

$F$  is *paranormal* if it is normal and also preserves compositions of the form  $g_* \bullet v$

$$\phi(g_*, v) : F(g_*) \bullet F(v) \longrightarrow F(g_* \bullet v) \quad \text{iso}$$

## Theorem

If  $F$  is lax paranormal, then it extends to  $F^{ret} : \mathbb{A}^{ret} \longrightarrow \mathbb{B}^{ret}$ , oplax paranormal

## Back to duals

### Theorem

*If  $\mathbb{A}$  has companions and left duals, the left dual is a lax normal double functor which is the identity on objects and horizontal arrows*

$$\bullet ( ) : \mathbb{A}^{ret\ co} \longrightarrow \mathbb{A}$$

- The proof uses strong universality



# Functoriality of $\backslash$

A cell  $\alpha$  and a retrocell  $\beta$  as in

$$\begin{array}{ccc}
 A & \xrightarrow{f} & A' \\
 \downarrow z & \Rightarrow \alpha & \downarrow z' \\
 C & \xrightarrow{h} & C' \\
 \uparrow y & \Leftarrow \beta & \uparrow y' \\
 B & \xrightarrow{g} & B'
 \end{array}$$

produce a cell

$$\begin{array}{ccc}
 A & \xrightarrow{f} & A' \\
 \downarrow y \backslash z & \Rightarrow \beta \backslash \alpha & \downarrow y' \backslash z' \\
 B & \xrightarrow{g} & B'
 \end{array}$$

given by

$$\begin{array}{ccccccc}
 A & \xlongequal{\quad} & A & \xlongequal{\quad} & A & \xrightarrow{f} & A' \\
 \downarrow y \backslash z & = & \downarrow y \backslash z & & \downarrow & & \downarrow \\
 B & \xlongequal{\quad} & B & \xrightarrow{\epsilon} & \bullet z & \xrightarrow{\alpha} & \bullet z' \\
 \downarrow g_* & & \downarrow y & & \downarrow & & \downarrow \\
 B' & \xrightarrow{\beta} & C & \xlongequal{\quad} & C & \xrightarrow{h} & C' \\
 \downarrow y' & & \downarrow h_* & = & \downarrow h_* & \lrcorner & \downarrow \text{id}_{C'} \\
 C' & \xlongequal{\quad} & C' & \xlongequal{\quad} & C' & \xlongequal{\quad} & C'
 \end{array}$$

## Theorem

$\backslash$  is functorial in both variables, covariant in the top variable and retrovariant in the bottom one